

Introduction to Intel x86-64 Assembly, Architecture, Applications, & Alliteration

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Control Flow

- Two forms of control flow
 - Conditional - go somewhere if a condition is met. Think “if”s, switches, loops
 - Unconditional - go somewhere no matter what. Function calls, goto, exceptions, interrupts.
- We’ve already seen procedure calls manifest themselves as call/ret, let’s see how goto manifests itself in asm.

GotoExample.c

```
//Goto example
#include <stdio.h>
int main(){
    goto mylabel;
    printf("skipped\n");
mylabel:
    printf("goto ftw!\n");
    return 0xf00d;
}
```

main:

00000000140001010	sub	rsp,28h
★ 00000000140001014	jmp	00000000140001023
00000000140001016	lea	rcx,[40006000h]
0000000014000101D	call	qword ptr [40008368h]
\$mylabel:		
00000000140001023	lea	rcx,[40006010h]
0000000014000102A	call	qword ptr [40008368h]
00000000140001030	mov	eax,0F00Dh
00000000140001035	add	rsp,28h
00000000140001039	ret	

Visual Studio “RIP-relative” display errors.
Goto “RIPRelativeAddressing” slides



JMP - Jump

- Change rip to the given address
- Main forms of the address
 - Short relative (1 byte displacement from end of the instruction)
 - “jmp 0000000140001023” doesn’t have the number 0000000140001023 anywhere in it, it’s really “jmp 0x0E bytes forward”
 - Some disassemblers will indicate this with a mnemonic by writing it as “jmp short”
 - Near relative (4 byte displacement from current eip)
 - Absolute (hardcoded address in instruction)
 - Absolute Indirect (address calculated with r/m32)
 - TODO CONFIRM R/M64
- jmp -2 == infinite loop for short relative jmp :)

GotoExample.c takeaways

- goto == jmp in asm :)

```
//Goto example
#include <stdio.h>
int main(){
    goto mylabel;
    printf("skipped\n");
mylabel:
    printf("goto ftw!\n");
    return 0xf00d;
}

main:
00000000140001010 sub    rsp,28h
00000000140001014 jmp     00000000140001023
00000000140001016 lea     rcx,[40006000h]
0000000014000101D call    qword ptr [40008368h]
$mylabel:
00000000140001023 lea     rcx,[40006010h]
0000000014000102A call    qword ptr [40008368h]
00000000140001030 mov     eax,0F00Dh
00000000140001035 add     rsp,28h
00000000140001039 ret
```

IfExample.c

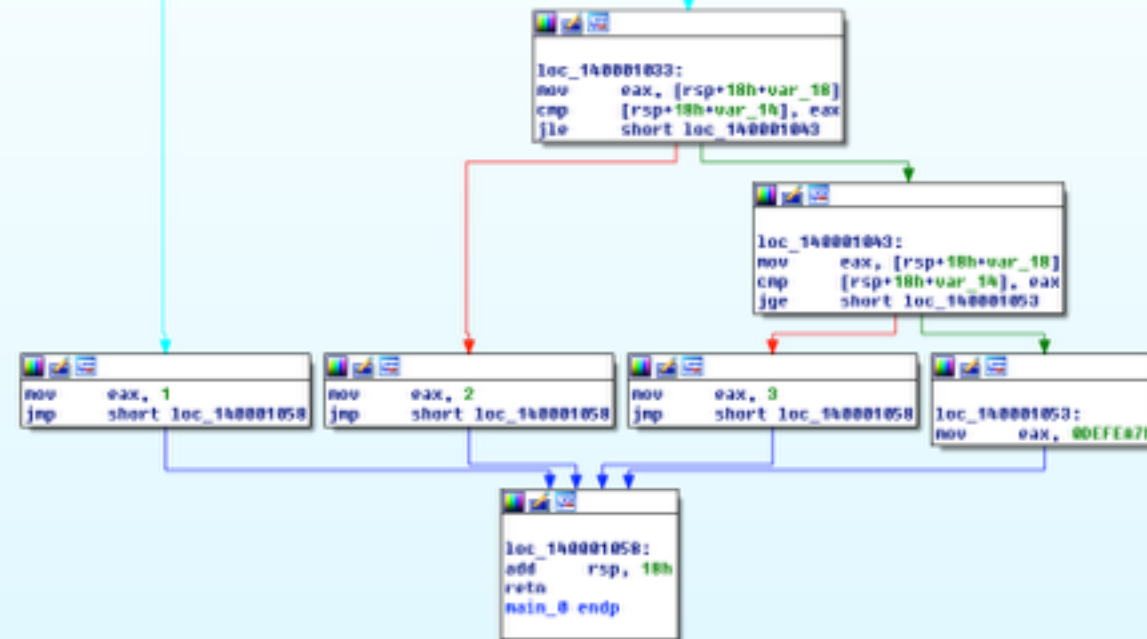
```
int main(){
    int a=1, b=2;
    if(a == b){
        return 1;
    }
    if(a > b){
        return 2;
    }
    if(a < b){
        return 3;
    }
    return 0xdefea7;
}
```

	main:	00000000140001010	sub	rsp,18h
		00000000140001014	mov	dword ptr [rsp+4],1
		0000000014000101C	mov	dword ptr [rsp],2
		00000000140001023	mov	eax,dword ptr [rsp]
		00000000140001026	cmp	dword ptr [rsp+4],eax
		0000000014000102A	jne	00000000140001033
		0000000014000102C	mov	eax,1
		00000000140001031	jmp	00000000140001058
		00000000140001033	mov	eax,dword ptr [rsp]
		00000000140001036	cmp	dword ptr [rsp+4],eax
		0000000014000103A	jle	00000000140001043
		0000000014000103C	mov	eax,2
		00000000140001041	jmp	00000000140001058
		00000000140001043	mov	eax,dword ptr [rsp]
		00000000140001046	cmp	dword ptr [rsp+4],eax
		0000000014000104A	jge	00000000140001053
		0000000014000104C	mov	eax,3
		00000000140001051	jmp	00000000140001058
		00000000140001053	mov	eax,0DEFEA7h
		00000000140001058	add	rsp,18h
		0000000014000105C	ret	

Jcc {

```
main_0 proc near
var_18= dword ptr -18h
var_14= dword ptr -14h
sub    rsp, 18h
mov     [rsp+18h+var_14], 1
mov     [rsp+18h+var_18], 2 ;
mov     eax, [rsp+18h+var_18]
cmp     [rsp+18h+var_14], eax
jnz     short loc_140001033
```

Ghost of Xmas Future:
Tools you won't get to use today
generate a Control Flow Graph (CFG)
which looks much nicer.
Not that that helps you. Just sayin' :)





Jcc - Jump If Condition Is Met

- There are more than 4 pages of conditional jump types! Luckily a bunch of them are synonyms for each other.
- JNE == JNZ (Jump if not equal, Jump if not zero, both check if the Zero Flag (ZF) == 0)

Some Notable Jcc Instructions

- JZ/JE: if ZF == 1
- JNZ/JNE: if ZF == 0
- JLE/JNG : if ZF == 1 or SF != OF
- JGE/JNL : if SF == OF
- JBE/JNA: if CF == 1 OR ZF == 1
- JB: if CF == 1
- Note: Don't get hung up on memorizing which flags are set for what. More often than not, you will be running code in a debugger, not just reading it. In the debugger you can just look at rflags and/or watch whether it takes a jump.

Some Notable Jcc Instructions

- Mnemonic translations
- A = above, unsigned notion
- B = below, unsigned notion
- G = greater than, signed notion
- L = less than, signed notion
- E = Equal (same as Z, zero flag set)
- N = Not (like “Not less than:” JNL)

Some Notable Jcc Instructions

- Mnemonic translations
- **B** = below, unsigned notion
- **A** = above, unsigned notion
- **N** = Not (like “Not less than:” JNL)
- **G** = greater than, signed notion
- **L** = less than, signed notion
- **E** = Equal (same a Z, zero flag set)



Final Fantasy 7
Crystal Bangle
Unlocked!

Flag setting

- Before you can do a conditional jump, you need something to set the condition flags for you.
- Typically done with CMP, TEST, or whatever instructions are already inline and happen to have flag-setting side-effects



CMP - Compare Two Operands

- “The comparison is performed by subtracting the second operand from the first operand and then setting the status flags in the same manner as the SUB instruction.”
- What’s the difference from just doing SUB? Difference is that with SUB the result has to be stored somewhere. With CMP the result is computed, the flags are set, but the result is discarded. Thus this only sets flags and doesn’t mess up any of your registers.
- Modifies CF, OF, SF, ZF, AF, and PF
- (as we already saw, SUB modifies all those too)

IfExample.c takeaways

- Conditional logic, like if statements, manifests in assembly as conditional jumps (Jcc). “If condition true, jump there, else fall through”
- Conditions involving in/equality are often checked with a CMP instruction, which is the same thing as a SUB, but it just throws the results away after the relevant RFLAGS bits are set
- The RFLAGS bits are fundamentally what are checked by the Jccs

```
int main(){
    int a=1, b=2;
    if(a == b){
        return 1;
    }
    if(a > b){
        return 2;
    }
    if(a < b){
        return 3;
    }
    return 0xdefea7;
}

main:
00000000140001010 sub    rsp,18h
00000000140001014 mov     dword ptr [rsp+4],1
0000000014000101C mov     dword ptr [rsp],2
00000000140001023 mov     eax,dword ptr [rsp]
00000000140001026 cmp     dword ptr [rsp+4],eax
0000000014000102A jne     00000000140001033
0000000014000102C mov     eax,1
00000000140001031 jmp     00000000140001058
00000000140001033 mov     eax,dword ptr [rsp]
00000000140001036 cmp     dword ptr [rsp+4],eax
0000000014000103A jle     00000000140001043
0000000014000103C mov     eax,2
00000000140001041 jmp     00000000140001058
00000000140001043 mov     eax,dword ptr [rsp]
00000000140001046 cmp     dword ptr [rsp+4],eax
0000000014000104A jge     00000000140001053
0000000014000104C mov     eax,3
00000000140001051 jmp     00000000140001058
00000000140001053 mov     eax,0DEFEA7h
00000000140001058 add     rsp,18h
0000000014000105C ret
```

BitmaskExample.c

```
#define MASK 0x100

int main(){
    int a=0x1301;
    if(a & MASK){
        return 1;
    }
    else{
        return 2;
    }
}

main:
00000000140001010 sub    rsp,18h
00000000140001014 mov     dword ptr [rsp],1301h
0000000014000101B mov     eax,dword ptr [rsp]
0000000014000101E and     eax,100h
00000000140001023 test    eax,eax
00000000140001025 je      00000000140001030
00000000140001027 mov     eax,1
0000000014000102C jmp     00000000140001035
0000000014000102E jmp     00000000140001035
00000000140001030 mov     eax,2
00000000140001035 add     rsp,18h
00000000140001039 ret
```




TEST - Logical Compare

- “Computes the bit-wise AND of first operand (source 1 operand) and the second operand (source 2 operand) and sets the SF, ZF, and PF status flags according to the result.”
- Like CMP - sets flags, and throws away the result

BitmaskExample.c takeaways

- Conditions depending on bit tests (which is often expressed with boolean logic instructions) will often see the RFLAGS set with the CMP instruction. CMP is like AND, but throws the results away
- The reason for the extraneous jmp here is because it's unoptimized code so it's following a simpler set of asm construction rules

#define MASK 0x100	main:	0000000140001010	sub	rsp,18h
		0000000140001014	mov	dword ptr [rsp],1301h
int main(){		000000014000101B	mov	eax,dword ptr [rsp]
int a=0x1301;		000000014000101E	and	eax,100h
if(a & MASK){		0000000140001023	test	eax,eax
return 1;		0000000140001025	je	0000000140001030
}		0000000140001027	mov	eax,1
else{		000000014000102C	jmp	0000000140001035
return 2;		000000014000102E	jmp	0000000140001035
}		0000000140001030	mov	eax,2
}		0000000140001035	add	rsp,18h
		0000000140001039	ret	

Instructions we now know (17)

- NOP
- PUSH/POP
- CALL/RET
- MOV
- ADD/SUB
- IMUL
- MOVZX/MOVSX
- LEA
- JMP/Jcc (family)
- CMP/TEST